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C. Perkins
 Nokia Research Center
 E. Belding-Royer
 University of California, Santa Barbara
 S. Das
 University of Cincinnati
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Ad hoc On-Demand Distance Vector (AODV) Routing

Status of this Memo

This memo defines an Experimental Protocol for the Internet community. It does not specify an Internet standard of any kind. Discussion and suggestions for improvement are requested. Distribution of this memo is unlimited.

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Abstract

The Ad hoc On-Demand Distance Vector (AODV) routing protocol is intended for use by mobile nodes in an ad hoc network. It offers quick adaptation to dynamic link conditions, low processing and memory overhead, low network utilization, and determines unicast routes to destinations within the ad hoc network. It uses destination sequence numbers to ensure loop freedom at all times (even in the face of anomalous delivery of routing control messages), avoiding problems (such as "counting to infinity") associated with classical distance vector protocols.

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1. Introduction

The Ad hoc On-Demand Distance Vector (AODV) algorithm enables dynamic, self-starting, multihop routing between participating mobile nodes wishing to establish and maintain an ad hoc network. AODV allows mobile nodes to obtain routes quickly for new destinations, and does not require nodes to maintain routes to destinations that are not in active communication. AODV allows mobile nodes to respond to link breakages and changes in network topology in a timely manner. The operation of AODV is loop-free, and by avoiding the Bellman-Ford "counting to infinity" problem offers quick convergence when the ad hoc network topology changes (typically, when a node moves in the network). When links break, AODV causes the affected set of nodes to be notified so that they are able to invalidate the routes using the lost link.

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One distinguishing feature of AODV is its use of a destination sequence number for each route entry. The destination sequence number is created by the destination to be included along with any route information it sends to requesting nodes. Using destination

sequence numbers ensures loop freedom and is simple to program. Given the choice between two routes to a destination, a requesting node is required to select the one with the greatest sequence number.

2. Overview

Route Requests (RREQs), Route Replies (RREPs), and Route Errors (RERRs) are the message types defined by AODV. These message types are received via UDP, and normal IP header processing applies. So, for instance, the requesting node is expected to use its IP address as the Originator IP address for the messages. For broadcast messages, the IP limited broadcast address (255.255.255.255) is used. This means that such messages are not blindly forwarded. However, AODV operation does require certain messages (e.g., RREQ) to be disseminated widely, perhaps throughout the ad hoc network. The range of dissemination of such RREQs is indicated by the TTL in the IP header. Fragmentation is typically not required.

As long as the endpoints of a communication connection have valid routes to each other, AODV does not play any role. When a route to a new destination is needed, the node broadcasts a RREQ to find a route to the destination. A route can be determined when the RREQ reaches either the destination itself, or an intermediate node with a 'fresh enough' route to the destination. A 'fresh enough' route is a valid route entry for the destination whose associated sequence number is at least as great as that contained in the RREQ. The route is made available by unicasting a RREP back to the origination of the RREQ. Each node receiving the request caches a route back to the originator of the request, so that the RREP can be unicast from the destination along a path to that originator, or likewise from any intermediate node that is able to satisfy the request.

Nodes monitor the link status of next hops in active routes. When a link break in an active route is detected, a RERR message is used to notify other nodes that the loss of that link has occurred. The RERR message indicates those destinations (possibly subnets) which are no longer reachable by way of the broken link. In order to enable this reporting mechanism, each node keeps a "precursor list", containing the IP address for each its neighbors that are likely to use it as a next hop towards each destination. The information in the precursor lists is most easily acquired during the processing for generation of a RREP message, which by definition has to be sent to a node in a precursor list (see section 6.6). If the RREP has a nonzero prefix

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length, then the originator of the RREQ which solicited the RREP information is included among the precursors for the subnet route (not specifically for the particular destination).

A RREQ may also be received for a multicast IP address. In this document, full processing for such messages is not specified. For example, the originator of such a RREQ for a multicast IP address may have to follow special rules. However, it is important to enable

correct multicast operation by intermediate nodes that are not enabled as originating or destination nodes for IP multicast addresses, and likewise are not equipped for any special multicast protocol processing. For such multicast-unaware nodes, processing for a multicast IP address as a destination IP address MUST be carried out in the same way as for any other destination IP address.

AODV is a routing protocol, and it deals with route table management. Route table information must be kept even for short-lived routes, such as are created to temporarily store reverse paths towards nodes originating RREQs. AODV uses the following fields with each route table entry:

- Destination IP Address
- Destination Sequence Number
- Valid Destination Sequence Number flag
- Other state and routing flags (e.g., valid, invalid, repairable, being repaired)
- Network Interface
- Hop Count (number of hops needed to reach destination)
- Next Hop
- List of Precursors (described in Section 6.2)
- Lifetime (expiration or deletion time of the route)

Managing the sequence number is crucial to avoiding routing loops, even when links break and a node is no longer reachable to supply its own information about its sequence number. A destination becomes unreachable when a link breaks or is deactivated. When these conditions occur, the route is invalidated by operations involving the sequence number and marking the route table entry state as invalid. See section 6.1 for details.

3. AODV Terminology

This protocol specification uses conventional meanings [1] for capitalized words such as MUST, SHOULD, etc., to indicate requirement levels for various protocol features. This section defines other terminology used with AODV that is not already defined in [3].

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active route

A route towards a destination that has a routing table entry that is marked as valid. Only active routes can be used to forward data packets.

broadcast

Broadcasting means transmitting to the IP Limited Broadcast address, 255.255.255.255. A broadcast packet may not be blindly forwarded, but broadcasting is useful to enable dissemination of AODV messages throughout the ad hoc network.

destination

An IP address to which data packets are to be transmitted. Same as "destination node". A node knows it is the destination node for a typical data packet when its address appears in the appropriate field of the IP header. Routes for destination nodes are supplied by action of the AODV protocol, which carries the IP address of the desired destination node in route discovery messages.

forwarding node

A node that agrees to forward packets destined for another node, by retransmitting them to a next hop that is closer to the unicast destination along a path that has been set up using routing control messages.

forward route

A route set up to send data packets from a node originating a Route Discovery operation towards its desired destination.

invalid route

A route that has expired, denoted by a state of invalid in the routing table entry. An invalid route is used to store previously valid route information for an extended period of time. An invalid route cannot be used to forward data packets, but it can provide information useful for route repairs, and also for future RREQ messages.

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originating node

A node that initiates an AODV route discovery message to be processed and possibly retransmitted by other nodes in the ad hoc network. For instance, the node initiating a Route Discovery process and broadcasting the RREQ message is called the originating node of the RREQ message.

reverse route

A route set up to forward a reply (RREP) packet back to the originator from the destination or from an intermediate node having a route to the destination.

sequence number

A monotonically increasing number maintained by each originating node. In AODV routing protocol messages, it is used by other nodes to determine the freshness of the information contained from the originating node.

valid route

See active route.

4. Applicability Statement

The AODV routing protocol is designed for mobile ad hoc networks with populations of tens to thousands of mobile nodes. AODV can handle low, moderate, and relatively high mobility rates, as well as a variety of data traffic levels. AODV is designed for use in networks where the nodes can all trust each other, either by use of preconfigured keys, or because it is known that there are no malicious intruder nodes. AODV has been designed to reduce the dissemination of control traffic and eliminate overhead on data traffic, in order to improve scalability and performance.

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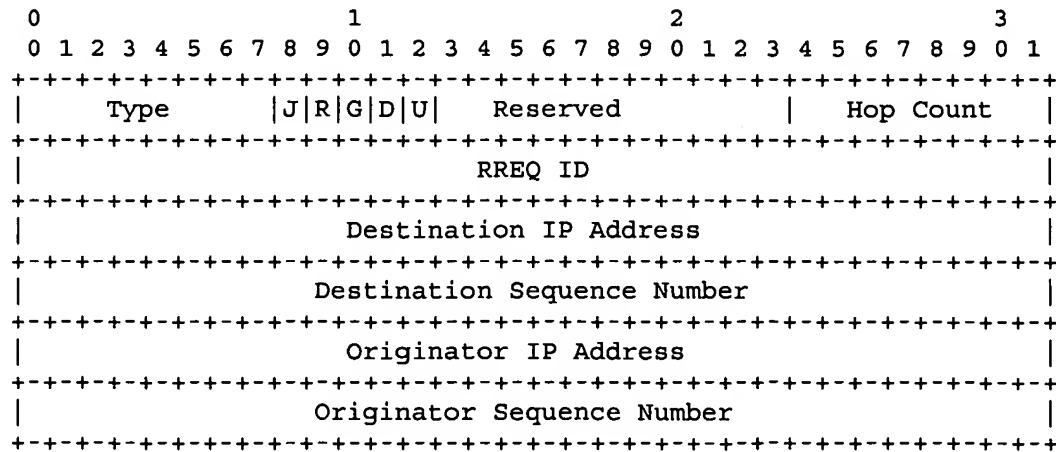
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5. Message Formats

5.1. Route Request (RREQ) Message Format



The format of the Route Request message is illustrated above, and contains the following fields:

Type	1
J	Join flag; reserved for multicast.
R	Repair flag; reserved for multicast.
G	Gratuitous RREP flag; indicates whether a gratuitous RREP should be unicast to the node specified in the Destination IP Address field (see sections 6.3, 6.6.3).
D	Destination only flag; indicates only the destination may respond to this RREQ (see section 6.5).
U	Unknown sequence number; indicates the destination sequence number is unknown (see section 6.3).
Reserved	Sent as 0; ignored on reception.
Hop Count	The number of hops from the Originator IP Address to the node handling the request.

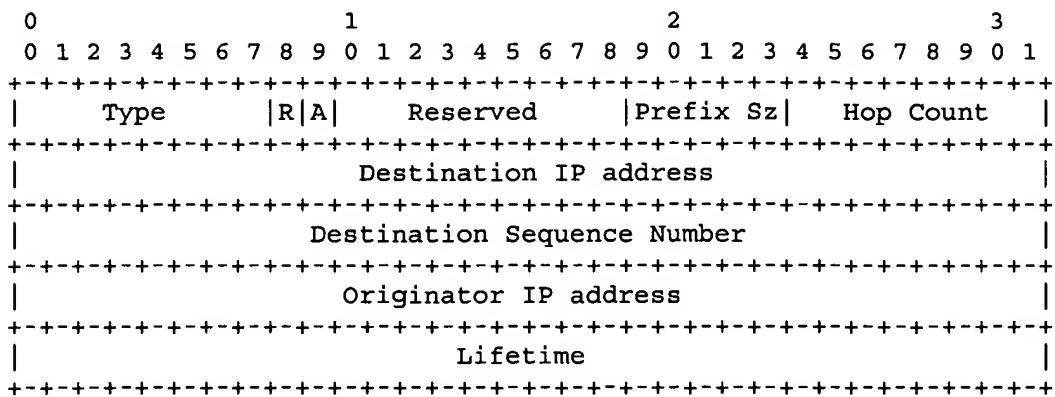
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RREQ ID	A sequence number uniquely identifying the particular RREQ when taken in conjunction with the originating node's IP address.
Destination IP Address	The IP address of the destination for which a route is desired.
Destination Sequence Number	The latest sequence number received in the past by the originator for any route towards the destination.
Originator IP Address	The IP address of the node which originated the Route Request.
Originator Sequence Number	The current sequence number to be used in the route entry pointing towards the originator of the route request.

5.2. Route Reply (RREP) Message Format



The format of the Route Reply message is illustrated above, and contains the following fields:

Type	2
R	Repair flag; used for multicast.
A	Acknowledgment required; see sections 5.4 and 6.7.
Reserved	Sent as 0; ignored on reception.

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Prefix Size If nonzero, the 5-bit Prefix Size specifies that the indicated next hop may be used for any nodes with the same routing prefix (as defined by the Prefix Size) as the requested destination.

Hop Count The number of hops from the Originator IP Address to the Destination IP Address. For multicast route requests this indicates the number of hops to the multicast tree member sending the RREP.

Destination IP Address
The IP address of the destination for which a route is supplied.

Destination Sequence Number
The destination sequence number associated to the route.

Originator IP Address
The IP address of the node which originated the RREQ for which the route is supplied.

Lifetime The time in milliseconds for which nodes receiving the RREP consider the route to be valid.

Note that the Prefix Size allows a subnet router to supply a route for every host in the subnet defined by the routing prefix, which is determined by the IP address of the subnet router and the Prefix

Size. In order to make use of this feature, the subnet router has to guarantee reachability to all the hosts sharing the indicated subnet prefix. See section 7 for details. When the prefix size is nonzero, any routing information (and precursor data) MUST be kept with respect to the subnet route, not the individual destination IP address on that subnet.

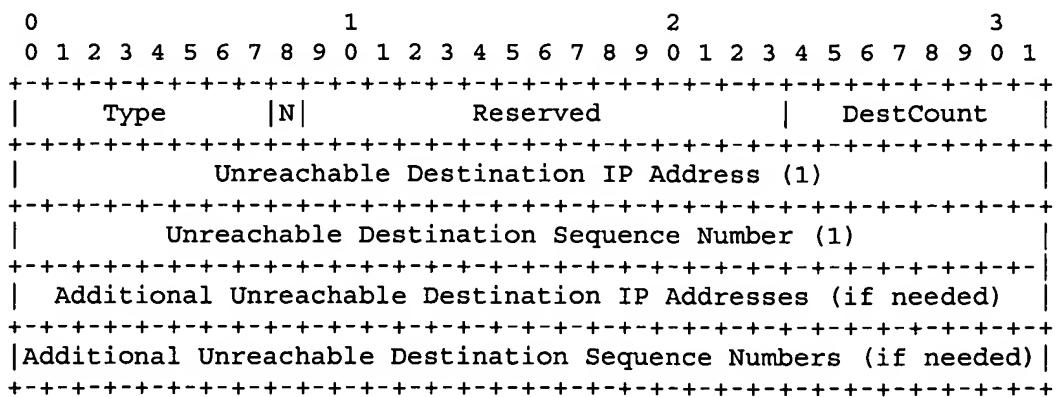
The 'A' bit is used when the link over which the RREP message is sent may be unreliable or unidirectional. When the RREP message contains the 'A' bit set, the receiver of the RREP is expected to return a RREP-ACK message. See section 6.8.

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5.3. Route Error (RERR) Message Format



The format of the Route Error message is illustrated above, and contains the following fields:

Type	3
N	No delete flag; set when a node has performed a local repair of a link, and upstream nodes should not delete the route.
Reserved	Sent as 0; ignored on reception.
DestCount	The number of unreachable destinations included in the message; MUST be at least 1.
Unreachable	Destination IP Address The IP address of the destination that has become

unreachable due to a link break.

Unreachable Destination Sequence Number

The sequence number in the route table entry for the destination listed in the previous Unreachable Destination IP Address field.

The RERR message is sent whenever a link break causes one or more destinations to become unreachable from some of the node's neighbors. See section 6.2 for information about how to maintain the appropriate records for this determination, and section 6.11 for specification about how to create the list of destinations.

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5.4. Route Reply Acknowledgment (RREP-ACK) Message Format

The Route Reply Acknowledgment (RREP-ACK) message MUST be sent in response to a RREP message with the 'A' bit set (see section 5.2). This is typically done when there is danger of unidirectional links preventing the completion of a Route Discovery cycle (see section 6.8).

0		1
0 1 2 3 4 5 6 7 8 9 0	1 2 3 4 5	
+-+-+-+---+---+---+---+---+---+		
Type	Reserved	
+-+-+-+---+---+---+---+---+---+		

Type 4

Reserved Sent as 0; ignored on reception.

6. AODV Operation

This section describes the scenarios under which nodes generate Route Request (RREQ), Route Reply (RREP) and Route Error (RERR) messages for unicast communication towards a destination, and how the message data are handled. In order to process the messages correctly, certain state information has to be maintained in the route table entries for the destinations of interest.

All AODV messages are sent to port 654 using UDP.

6.1. Maintaining Sequence Numbers

Every route table entry at every node MUST include the latest information available about the sequence number for the IP address of the destination node for which the route table entry is maintained. This sequence number is called the "destination sequence number". It is updated whenever a node receives new (i.e., not stale) information